Generating high-performance custom floating-point pipelines
Florent de Dinechin, Cristian Klein, Bogdan Pasca

To cite this version:
Florent de Dinechin, Cristian Klein, Bogdan Pasca. Generating high-performance custom floating-point pipelines. Field Programmable Logic and Applications, Aug 2009, Prague, Czech Republic. ensl-00379154v2

HAL Id: ensl-00379154
https://hal-ens-lyon.archives-ouvertes.fr/ensl-00379154v2
Submitted on 28 Apr 2009

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L’archive ouverte pluridisciplinaire HAL, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d’enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.
ABSTRACT

Custom operators, working at custom precisions, are a key ingredient to fully exploit the FPGA flexibility advantage for high-performance computing. Unfortunately, such operators are costly to design, and application designers tend to rely on less efficient off-the-shelf operators. To address this issue, an open-source architecture generator framework is introduced. Its salient features are an easy learning curve from VHDL, the ability to embed arbitrary synthesizable VHDL code, portability to mainstream FPGA targets from Xilinx and Altera, automatic management of complex pipelines with support for frequency-directed pipeline, automatic test-bench generation. This generator is presented around the simple example of a collision detector, which it significantly improves in accuracy, DSP count, logic usage, frequency and latency with respect to an implementation using standard floating-point operators.

1. INTRODUCTION

1.1. Flexible operators for FPGAs

FPGA-based coprocessors are available from a variety of vendors and may be used for accelerating intensive computations, including floating-point (FP) ones. On matrix multiplication, their floating-point performance barely surpasses that of a contemporary processor [1], using tens of operators on the FPGA to compensate their much slower frequency. Besides, FPGAs are in competitions with GPUs here. There is much more acceleration potential in operations for which processors or GPU have no dedicated hardware, for instance double-precision elementary functions [2] or special functions used in the generation of floating-point random numbers [3]. Besides, operators can also be specialized in FPGAs. For example, a squarer requires less logic than a multiplier; a floating-point multiplication by the constant 2.0 boils down to adding one to the exponent (a 12-bit addition for double-precision), and need not use a full-blown FP multiplier as it does in a processor. Indeed it is possible to build an optimized architecture for any multiplication by a constant [4]. Finally, and building upon the previous optimizations, operators can be fused on an FPGA [5]. The running example illustrating this article is such a fused operator, a collision detection predicate, testing if the norm of a 3D vector is larger than some threshold.

The recipe for implementing such operators is to respect their floating-point interface, but perform most internal computations in fixed point for which FPGAs are highly optimized. This recipe is also used in block-floating floating-point approaches [6] and in recent experimental C-to-FPGA datapath synthesis tools [5].

There are many more opportunities for floating-point on FPGAs [7]. The goal of the FloPoCo project\footnote{www.ens-lyon.fr/LIP/arenaire/Ware/FloPoCo/} is to study and develop such FPGA-specific Floating-Point Cores. FloPoCo is both a software framework and a generator of arithmetic cores using this framework. It currently offers about 20 operators, from simple ones such as shifters or integer adders to very complex ones such as floating-point exp and log. Some of the original operators have been described elsewhere [8, 2, 4, 9], and this article presents the framework itself. FloPoCo is distributed under the LGPL, and interested readers are welcome to try it, use it and improve it.

1.2. The arithmetic context

Our definition of an arithmetic operator is much broader than the usual +, −, × and /. In this work, an arithmetic operator is the implementation of any useful mathematical function, be it ln(− ln(x)) [3] or $\frac{x}{\sqrt{x^2+y^2+z^2}}$.

Although very general, this mathematical definition has
many important consequences when designing an operator. To start with, being the implementation of a function, an operator usually involves no feedback loop or state machine. This restriction enables a simple framework for quickly developing high-performance pipelined operators out of combinatorial ones. This is a major contribution of this work and will be exposed in Section 3.2.

When an operator is defined as the implementation of a mathematical function, the error analysis and design-space exploration may be much more accurate than when it is the implementation of a C program. Finally, the mathematical function also defines a reference for testing, as will be detailed in Section 4.

1.3. From libraries to generators

Although early FP operators were proposed as VHDL or Verilog libraries, the current trend is to shift to generators of operators (see [10] and references therein). A generator is a program that inputs user specifications, performs any relevant architectural exploration and construction (sometimes down to pre-placement), and outputs the architecture in a synthesizable format. Most FPGA designers are familiar with the Xilinx core generator tool, which to our knowledge pioneered this approach, or its Altera MegaWizard counterpart. A generator may simply wrap a library – and for the simplest operators there is no need for more – but it can also be much more powerful.

Firstly, the size of a library including multipliers by all the possible constants would be infinite, but the generation of an architecture for such a multiplier may be automated as a program that inputs the constant [4]. The same holds for arbitrary function generators [8].

Secondly, generators allow for greater parameterization and flexibility. Whether the best operator is a slow and small one or a faster but larger one depends on the context. FPGAs also allow flexibility in precision: arithmetic cores should be parameterized by the bit-widths of their inputs and outputs. We are also concerned about optimizing for different hardware targets, with different LUT structure, memory and DSP features, etc. The more flexible a generator, the more future-proof.

Lastly, generators may perform arbitrarily complex design-space exploration, automated error analysis, and optimization [8, 2].

1.4. Design choices for FloPoCo

A generator could in principle be written in VHDL or Verilog, but these languages are not ideally suited for that. In particular, complex programs turn out to be inefficient and difficult to debug using the simulation/synthesis paradigm of hardware description languages. Still, an architecture generator needs a back-end to actually implement the resulting circuit. The most elegant solution is to write a generator as an overlay on a software-based HDL such as SystemC, JBits, HandelC or JHDL (among many others). The advantages are a preexisting abstraction of a circuit, and simple integration with a one-step compilation process. The inconvenience is that most of these languages are still relatively confidential and restricted in the FPGAs they support. Even SystemC synthesizers are far from being commonplace yet. Basing our generator on a vendor generator would be an option, but would mean restricting it to one FPGA family.

The proposed generator took a less elegant, but more universal route. The generator is written in a mainstream programming language (we chose C++, mostly for compatibility with other previous work), and it outputs operators in a mainstream HDL (we chose standard synthesizable VHDL). We thus get a portable generator, and the generated operators can be integrated into most projects, simulated using mainstream simulators, and synthesized for any FPGA using the vendor back-end tools. Section 3.3 will discuss how they can nevertheless be optimized to a given FPGA target.

2. A SIMPLE MOTIVATING EXAMPLE

This framework has been used to write and pipeline very large components, such as the floating-point logarithm described in [2]. We choose to discuss here a simpler, but still representative example: the implementation of a collision detection predicate, testing if the norm of a 3D vector is larger than some threshold. Let \( x, y, z \) and \( r \) be four floating-point inputs, and let us design an operator testing whether \( x^2 + y^2 + z^2 < r^2 \).

A first option is to assemble three FP multipliers and two FP adders. The final comparison can be performed by an integer subtraction on the concatenation of the exponent and the mantissa – a well-known property of IEEE-754 FP formats is that positive FP numbers are ordered in the lexicographic order of their binary representation.

Let us now design a specific operator for this computation. There are many optimization opportunities with respect to the previous solution.

- As already mentioned, the squarers should be simpler than multipliers. They will in particular use less DSP blocks.
- We only add positive numbers. In an off-the-shelf floating-point adder, roughly half the logic is dedicated to effective subtraction, and can be saved here. An optimizing synthesizer would probably perform this optimization, but it will not, for instance, reduce
Overview of the custom floating-point collision operator.

- The significands of $x^2$, $y^2$ and $z^2$ are aligned before addition. This may be done in parallel, reducing pipeline depth and register usage with respect to an implementation using two floating-point adders in sequence.
- A lot of logic is dedicated to rounding intermediate results, and can be saved by considering the compound operator as an atomic one [5]. We obtain an operator that is not only simpler, but also more accurate.

Actually, it is common for a fused operator to be more accurate than the combination of FP ones. Here, the idea is to evaluate $x^2 + y^2 + z^2$ with 1ulp (unit in the last place) accuracy before the comparison, where a combination of standard operators leads to a worst-case error larger than two ulps. More subtly, with an operator built by assembling two floating-point adders, there are some rare cases when the value of the predicate will change when one exchanges $x$ and $z$. This is due to the two consecutive roundings. We will be able to design an operator which doesn’t have such asymmetries. It is even possible to design an exact predicate, that doesn’t perform any rounding before the comparison – it probably doesn’t make much sense unless $x$, $y$ and $z$ are computed exactly themselves.

The chosen architecture is depicted on Figure 1. For more details the interested reader is invited to have a look at the Collision operator in the FloPoCo distribution. This architecture computes the three squares, truncates them to $w_F + g$ bits, then aligns them to the largest one and adds them. Worst-case error analysis shows that there are 5 truncation errors in this computation (the three products, and bits discarded by the two shifters). The number of guard bits $g$ is therefore set to $g = \lceil \log_2(5) \rceil = 3$ so that the accumulated error is always smaller than the unit in the last place of $r^2$. Comparison is performed as integer subtraction as previously, except that $r^2$ is extended to $w_F + g$ bits.

Table 1 shows quantitatively the savings brought by this approach. All the metrics are significantly improved (from 25 to 89 %) by operator specialization and fusion. The (8,23) FP version implemented using LogiCore illustrates that FloPoCo provides FP operators of comparable quality.

Another meaningful result is design time. Assembling FP operators took two hours, for a pipeline that is fully parameterized by $w_E$, $w_F$ and target frequency. Assembling the Logicore version took even less time, but one doesn’t obtain a parameterized design. The optimized version took two days to write, including the testbench. It may seem a lot of design time, but most of the work concerns the sum-of-square operator, which will also be used for future coarser-grain operators such as $\sqrt{x^2 + y^2 + z^2}$ or $\frac{x}{\sqrt{x^2 + y^2 + z^2}}$. And again, this design is fully parameterized, pipelined, and automatically optimized for a given frequency.

Let us now present in more details the framework used to obtain these results.

3. THE FRAMEWORK

In the following, we assume basic knowledge of object-oriented concepts with the C++ terminology. Figure 2 provides a very simplified overview of the FloPoCo class hierarchy.

3.1. Operators

The core class is Operator. From the circuit point of view, an Operator corresponds to a VHDL entity, but again, with restrictions and extensions specific to the arithmetic context. All the operators extend this class, including Collision,

Fig. 1. Overview of the custom floating-point collision operator.

Fig. 2. FloPoCo class hierarchy (very simplified overview)
but also some of its sub-components (shifters, squarers and adders) seen on Figure 1.

The main method of Operator is outputVHDL(), which prints out the VHDL code of an operator. To implement this virtual method for an operator, one may simply embed some existing VHDL code in the C++.

However, the standard way of using the framework is to rely on the default implementation of outputVHDL() provided by the Operator class. This default method takes the VHDL code of the operator architecture from the vhdl stream, an attribute of Operator. Thus, one designs the architecture by having the constructor outputting arbitrary VHDL code to vhdl. In addition, signals may be wrapped in methods from Operator. A first method is declare(), that handles signal declaration. Consider the following code that computes the difference of two exponents of size \( w_E \), in order to determine in XltY which is smaller.

\[
\text{vhdl} \text{<< declare("DEXY", wE+1) << " <= DEXY("<< wE<< ")- ("<< 0 << " & " << 0 << ") - ("<< 0 << " & " << EY << ")" << endl;} \\
\text{vhdl} \text{<< declare("XltY") << " <= " DEXY("<< wE<< ")" << endl;} \\
\]

Here the declare() method adds the signal to the signalList of the operator (so that outputVHDL() will declare it in the VHDL), then simply returns its first argument. Thus, called with \( w_E = 8 \), the code above simply puts in the vhdl stream the string:

\[
\text{DEXY} \leftarrow \left( \lfloor 0' \rfloor \text{ & EX} \right) - \left( \lfloor 0' \rfloor \text{ & EY} \right); \\
\text{XltY} \leftarrow \text{DEXY}(8); \\
\]

So far we have mostly obsfucated the VHDL with C++ constructs. Let us see how this enables easy pipelining.

### 3.2. Automatic pipeline management

We now consider the end of the normalise/pack box of Figure 1, which reassembles the result from the exponent datapath and the fraction datapath into a single bit vector. The exponent datapath is in general shorter (computed in less cycles), so there is a need for synchronisation. The cycle difference depends on the pipeline depths of the shifter and add components, which in turn depends on operand sizes, target hardware, and target frequency. Building a working pipeline for a given set of parameters is conceptually simple: simply insert registers to delay the result of the shortest datapath, so as to synchronise it with the result of the longest one. FloPoCo allows to express exactly that in a generic way. Consider the following code excerpt:

\[
\text{(...)} \\
\text{// at some cycle} \\
\text{vhdl} \text{<< declare("finalExp", wE+1) << " <= " ...} \\
\text{(...)} \\
\text{// at some other cycle} \\
\text{vhdl} \text{<< declare("finalFraction", wF+g) << " <= " ...} \\
\text{(...)} \\
\text{// enter next cycle} \\
\text{nextCycle();} \\
\text{vhdl} \text{<< declare("finalSoP", wE+wF+g) << " <= " ...} \\
\text{use("finalExp") << "(wE-1 downto 0)"} \\
\text{use("finalFraction") << ";", "};} \\
\]

During the construction of the VHDL in the vhdl stream, an Operator maintains a currentCycle attribute, initially equal to zero. The main function of nextCycle() is to increment currentCycle.

Every signal declared through addInput or declare has a cycle attribute, which represents the cycle at which this signal is computed. It is 0 for the inputs. For signals declared through declare(), it is set to currentCycle at the time declare() is invoked.

Every signal also possesses an attribute lifeSpan which indicates how many cycles it will need to be delayed. This attribute is initialized to 0, then possibly increased by use() as we will see below. When the lifeSpan of a signal \( x \) is greater than zero, outputVHDL() will create lifeSpan new signals, named \( x_{d1}, x_{d2} \) and so on, and insert registers between them. In other words, \( x_{d2} \) will hold the value of \( x \) delayed by 2 cycles.

| Precision \((w_E, w_F)\) | Slow version | | Fast version | |
|--------------------------|--------------| |---------------|--------------|
| \((8, 23)\) LogiCore FP  | 940 slices, 12 DSP | 20 cycles @ 210 MHz | 1282 slices, 20 DSP | 43 cycles @ 353 MHz |
| \((8, 23)\) FP            | 456 slices, 9 DSP  | 10 cycles @ 319 MHz | 1188 slices, 12 DSP | 29 cycles @ 289 MHz |
| \((8, 23)\) custom        | 1268 slices, 12 DSP | 20 cycles @ 171 MHz | 1874 slices, 12 DSP | 37 cycles @ 302 MHz |
| \((9, 32)\) FP            | 629 slices, 9 DSP  | 10 cycles @ 368 MHz | 640 slices, 9 DSP  | 13 cycles @ 368 MHz |
| \((9, 32)\) custom        | 2868 slices, 27 DSP | 20 cycles @ 106 MHz | 1845 slices, 18 DSP | 46 cycles @ 276 MHz |
| \((11, 52)\) FP           | 1532 slices, 18 DSP | 13 cycles @ 362 MHz | 37 cycles @ 302 MHz | 1532 slices, 18 DSP | 43 cycles @ 368 MHz |
| \((11, 52)\) custom       | 20 cycles @ 368 MHz | 37 cycles @ 302 MHz | 1532 slices, 18 DSP | 46 cycles @ 276 MHz |

Table 1. Some post-synthesis results for the Collision operator (on Virtex-4 xc4vlx15-12-sf363 using ISE 10.1 or 9.1, whichever gave the best results). Note that the custom version is also the most accurate. The FloPoCo command-line used to obtain these results is flopoco -frequency=f Collision \( w_E \ w_F \ o \) where \( f \) is 200 for the slow version and 400 for the fast version, and \( o \) is 0 for the FP version and 1 for the optimized version.
Wrapping a signal in \texttt{use()} has the following effect. First, \texttt{use("X")} will compare \texttt{currentCycle} and the \texttt{cycle} declared for \texttt{X}, which we note \texttt{X.cycle}.

- If they are equal, \texttt{use("X")} will simply return the string "X".
- If \texttt{currentCycle < X.cycle}, \texttt{use("X")} will emit an error message reporting that \texttt{X} is being used before the cycle at which it was defined.
- If \texttt{currentCycle > X.cycle}, \texttt{use("X")} will delay signal \texttt{X} by \texttt{n = currentCycle - X.cycle} cycles. Technically, if \texttt{n = 5} for instance, \texttt{use("X")} just returns "X.d5*. It also updates \texttt{X.lifespan} to be at least equal to \texttt{5}.

Operator defines several other functions to manage \texttt{currentCycle}, such as \texttt{setCycle(int n)}, \texttt{setCycleFromSignal(string s)} which sets the \texttt{currentCycle} to the \texttt{cycle} of the signal whose name is given as an argument (going back if needed), and \texttt{syncCycleFromSignal(string s)} which may only advance \texttt{currentCycle}. The latter allows to synchronise several signals by setting \texttt{currentCycle} to the max of their \texttt{cycle}.

It should be noted that this scheme still allows for arbitrary synthesizable VHDL – only the signal names have to be wrapped by \texttt{use()}, \texttt{declare()}, etc. Also, this scheme gracefully degrades to a combinatorial operator – in this mode \texttt{declare()} always declares one single signal and \texttt{use()} always returns its argument unmodified.

Code written under this scheme will automatically adapt to random insertions and suppressions of synchronisation barriers. Typically, if one is unhappy with the performance, one may decide to break the critical path by inserting a synchronisation barrier in it. This may be as simple as inserting one may decide to break the critical path by inserting a synchronisation barrier in it. This may be as simple as inserting a single \texttt{nextCycle()} in the code. FloPoCo takes care of the rest.

Finally, it is also possible to wrap in conditional statements some of the \texttt{nextCycle()} calls, so that the pipeline adapts to the frequency, the operator generic parameters, etc. For instance, the \texttt{IntAdder} or \texttt{Shifters} classes of FloPoCo produce operators whose pipeline depth depends on the target frequency. By using these pre-existing subcomponents, our \texttt{Collision} operator behaves the same, as Table 1 shows.

Space limitation prevents presenting in details how subcomponents are managed in FloPoCo. In short, one has to create an object instance for the subcomponent, then there is a \texttt{inPortMap()} method that is conceptually similar to \texttt{use()}, \texttt{anOutPortMap()} method conceptually similar to \texttt{declare()}, and an \texttt{instance()} method that returns the VHDL of the instance, and performs the bookkeeping required for \texttt{outputVHDL()} to declare the corresponding component. We invite the interested reader to have a look at \texttt{Collision.cpp}.

### 3.3. Targets

The \texttt{Target} class abstracts the features of actual FPGA chips. Classes representing real FPGA chips extend this class (we currently have classes for two very different FPGAs, Xilinx \texttt{Virtex4} and Altera \texttt{StratixII}). The idea is to declare abstract methods in \texttt{Target}, which are implemented in its subclasses, so that the same generator code fits all the targets. Of course, it is also possible to have a conditional statement that runs completely different code depending on the target – this will be the case for instance for the \texttt{IntMultiplier} class that builds large multipliers, because DSP capabilities are too variable from one target to the other. A \texttt{Target} is given as argument to the constructor of any operator.

The methods provided by the \texttt{Target} class can be semantically split into two categories:

- **Architecture-related methods** provide information about the architecture of the FPGA and are used in architectural exploration. For instance, \texttt{lutInputs()} returns the number of inputs of the FPGA’s LUTs.

- **Delay-related methods** provide approximative informations about computation time. For example, \texttt{adderDelay(int n)} returns the delay of an \texttt{n-bit} addition. These methods are used for frequency-directed pipelining (here for \texttt{IntAdder}). Some of these methods have an architecture-related dual, for example \texttt{suggestAdderSize(double delay)} that returns the size of an adder that will have the required delay.

The difficulty here is to find the right abstraction level for \texttt{Target}. On one hand, we do not hope to provide an exhaustive and detailed description of all the existing – and future – FPGAs. On the other hand, we do not need to: Vendor tools are very good at fitting a design to a given target, and we should rely on them. The complexity of exploiting the details of the target should be left to the back-end tools. \texttt{Target} is still expected to evolve a lot. Other methods will be added when needed, for instance to model the internal memory resources. Also, area estimation methods will assist the design-space exploration of future operators. And of course, the question of target abstraction will remain an open one forever, as new FPGA features keep appearing and new operators present new problems.

### 4. TESTBENCH GENERATION

The underlying mathematical nature of an arithmetic operator can be usefully exploited to make testbench generation
easy. For instance, the output from a floating-point opera-
tor may be defined as a mathematical function of the inputs
(addition or exponential) composed with one of the rounding
functions defined in the well established IEEE-754 standard
[11]. Defining expected outputs this way is not only simpler
than mimicking the architecture, it also minimizes the pos-
sibility of making the same mistake in both the operator and
its test bench. In FloPoCo, one has to overload the Oper-
ator virtual method emulate(). Thanks to the bit-accurate
MPFR library [12], this typically takes about ten lines.

Overloading emulate() is enough for FloPoCo to be
able to create a generic test bench using random inputs.
However, function analysis also allows for better, more
operator-specific test-case generation. Let us just take two
examples.

- A double-precision exponential returns \( +\infty \) for all
  inputs larger than 710 and returns 0 for all inputs
  smaller than \(-746\). In other terms, the most inter-
esting test domain for this function is when the input
  exponent is between \(-10\) and 10, a fraction of the full
  double-precision exponent domain \((-1024 \text{ to } 1023)\).
  Generating random 64-bit integers and using them as
  floating-point inputs would mean testing mostly the
  overflow/underflow logic, which is a tiny part of the
  operator.

- In a floating-point adder, if the difference between the
  exponents of the two operands is large, the adder will
  simply return the biggest of the two, and again this is
  the most probable situation when taking two random
  operands. Here it is better to generate random cases
  where the two operands have close exponents.

Such cases are managed by overloading the Oper-
ator method buildRandomTestCases(). Finally,
buildStandardTestCases() allows to test corner cases
which random testing has little chance to find.

5. CONCLUSION

This article introduced an open-source software framework
for generating high-quality, highly parameterized, pipelined
and flexible operators for FPGAs. This framework evolved
(and still evolves) in a deliberately bottom-up and practical
way focussed on the needs of arithmetic cores. It is based
on embedding VHDL code in a simple C++ framework.
Thus it provides the full power of the VHDL language. It
adds to it facilities for automatic pipeline construction and
synchronization, target-based and frequency-based optimi-
sation, and test bench generation.

In its current state, this generator is already an invaluable
booster for the core focus of our research, which is to de-
velop new arithmetic operators. In the future, a wider range

6. REFERENCES

[1] D. Strenski, J. Simkins, R. Walke, and R. Wittig, “Revaluat-
ing FPGAs for 64-bit floating-point calculations,” HPC wire,
May 2008.

ware floating-point elementary function,” in 18th Symposium

Monte Carlo in finance,” in Field Programmable Logic and

ger and floating-point constant multipliers for FPGAs,” in
Application-specific Systems, Architectures and Processors.
IEEE, 2008, pp. 239–244.

gas,” in Field Programmable Logic and Applications, 2008,

[6] R. Andraka, “Hybrid floating point technique yields 1.2 gi-
 gasample per second 32 to 2048 point floating point FFT in
a single FPGA,” in High Performance Embedded Computing
Workshop, 2006.

an, “When FPGAs are better at floating-point than micro-

fast hardware function evaluation,” in Application-specific
328–333.

[9] F. de Dinechin, B. Pasca, O. Creț, and R. Tudoran,
“An FPGA-specific approach to floating-point accumulation
and sum-of-products,” in Field-Programmable Technologies.

[10] D. Bakalis, K. Adaos, D. Lymperopoulos, M. Bellos, H. Ver-
gos, G. Alexiou, and D. Nikolos, “A core generator for arith-
cmetic cores and testing structures with a network interface,”
Journal of Systems Architecture, vol. 52, no. 1, pp. 1 – 12,
2006.

Arithmetic (also IEC 60559), 1985.

[12] L. Fousse, G. Hanrot, V. Lefèvre, P. Pélissier, and P. Zimmer-
mann, “MPFR: A multiple-precision binary floating-point li-
brary with correct rounding,” ACM Transactions on Mathe-