

Topological Complexity of omega-Powers: Extended Abstract

Olivier Finkel, Dominique Lecomte

► **To cite this version:**

Olivier Finkel, Dominique Lecomte. Topological Complexity of omega-Powers: Extended Abstract. Dagstuhl Seminar on "Topological and Game-Theoretic Aspects of Infinite Computations" 29.06.08 - 04.07.08, P. Hertling, V. Selivanov, W. Thomas, W. W. Wadge, K. Wagner, Dagstuhl, Germany. ensl-00319447v2

HAL Id: ensl-00319447

<https://hal-ens-lyon.archives-ouvertes.fr/ensl-00319447v2>

Submitted on 10 Sep 2008

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Topological Complexity of ω -Powers : Extended Abstract

Olivier Finkel¹ and Dominique Lecomte²

¹ *Equipe Modèles de Calcul et Complexité
Laboratoire de l'Informatique du Parallélisme***
CNRS et Ecole Normale Supérieure de Lyon
46, Allée d'Italie 69364 Lyon Cedex 07, France.

`Olivier.Finkel@ens-lyon.fr`

² *Equipe d'Analyse Fonctionnelle*
Université Paris 6
4, place Jussieu, 75 252 Paris Cedex 05, France
`dominique.lecomte@upmc.fr`

Keywords: Infinite words; ω -languages; ω -powers; Cantor topology; topological complexity; Borel sets; Borel ranks; complete sets; Wadge hierarchy; Wadge degrees; effective descriptive set theory; hyperarithmetical hierarchy

1 Introduction

The operation $V \rightarrow V^\omega$ is a fundamental operation over finitary languages leading to ω -languages. It produces ω -powers, i.e. ω -languages in the form V^ω , where V is a finitary language. This operation appears in the characterization of the class REG_ω of ω -regular languages (respectively, of the class CF_ω of context free ω -languages) as the ω -Kleene closure of the family REG of regular finitary languages (respectively, of the family CF of context free finitary languages) [Sta97a].

Since the set Σ^ω of infinite words over a finite alphabet Σ can be equipped with the usual Cantor topology, the question of the topological complexity of ω -powers of finitary languages naturally arises and has been posed by Niwinski [Niw90], Simonnet [Sim92], and Staiger [Sta97a]. A first task is to study the position of ω -powers with regard to the Borel hierarchy (and beyond to the projective hierarchy) [Sta97a,PP04].

It is easy to see that the ω -power of a finitary language is always an analytic set because it is either the continuous image of a compact set $\{0, 1, \dots, n\}^\omega$ for $n \geq 0$ or of the Baire space ω^ω .

It has been recently proved, that for each integer $n \geq 1$, there exist some ω -powers of context free languages which are Π_n^0 -complete Borel sets, [Fin01], and that there exists a context free language L such that L^ω is analytic but not Borel, [Fin03]. Notice that amazingly the language L is very simple to describe and it is accepted by a simple 1-counter automaton.

** UMR 5668 - CNRS - ENS Lyon - UCB Lyon - INRIA
LIP Research Report RR 2008-27

The first author proved in [Fin04] that there exists a finitary language V such that V^ω is a Borel set of infinite rank. It was also proved in [DF07] that there is a context free language W such that W^ω is Borel above Δ_ω^0 .

We proved in [FL07] the following very surprising result which shows that ω -powers exhibit a great topological complexity: for each non-null countable ordinal ξ , there exist some Σ_ξ^0 -complete ω -powers, and some Π_ξ^0 -complete ω -powers.

We consider also the Wadge hierarchy which is a great refinement of the Borel hierarchy. We get many more Wadge degrees of ω -powers, showing that for each ordinal $\xi \geq 3$, there are uncountably many Wadge degrees of ω -powers of Borel rank $\xi + 1$.

We show also, using some tools of effective descriptive set theory, that the main result of [FL07] has some effective counterparts.

All the proofs of the results presented here may be found in the conference paper [FL07] or in the preprint [FL08] which contains also some additional results.

2 Topology

We first give some notations for finite or infinite words, assuming the reader to be familiar with the theory of formal languages and of ω -languages, see [Tho90,Sta97a,PP04]. Let Σ be a finite or countable alphabet whose elements are called letters. A non-empty finite word over Σ is a finite sequence of letters: $x = a_0.a_1.a_2 \dots a_n$ where $\forall i \in [0; n]$ $a_i \in \Sigma$. We shall denote $x(i) = a_i$ the $(i+1)^{th}$ letter of x . The length of x is $|x| = n+1$. The empty word has 0 letters. Its length is 0. The set of finite words over Σ is denoted $\Sigma^{<\omega}$. A (finitary) language L over Σ is a subset of $\Sigma^{<\omega}$. The usual concatenation product of u and v will be denoted by $u \frown v$ or just uv .

The first infinite ordinal is ω . An ω -word over Σ is an ω -sequence $a_0a_1 \dots a_n \dots$, where for all integers $i \geq 0$ $a_i \in \Sigma$. When σ is an ω -word over Σ , we write $\sigma = \sigma(0)\sigma(1) \dots \sigma(n) \dots$. The set of ω -words over the alphabet Σ is denoted by Σ^ω . An ω -language over an alphabet Σ is a subset of Σ^ω . The concatenation product is also extended to the product of a finite word u and an ω -word v : the infinite word $u.v$ or $u \frown v$ is then the ω -word such that: $(uv)(k) = u(k)$ if $k < |u|$, and $(u.v)(k) = v(k - |u|)$ if $k \geq |u|$.

The prefix relation is denoted \prec : the finite word u is a prefix of the finite word v (respectively, the infinite word v), denoted $u \prec v$, if and only if there exists a finite word w (respectively, an infinite word w), such that $v = u \frown w$.

For a finitary language $V \subseteq \Sigma^{<\omega}$, the ω -power of V is the ω -language

$$V^\omega = \{u_1 \dots u_n \dots \in \Sigma^\omega \mid \forall i \geq 1 \ u_i \in V\}$$

We recall now some notions of topology, assuming the reader to be familiar with basic notions which may be found in [Kur66,Mos80,Kec95,LT94,Sta97a,PP04].

There is a natural metric on the set Σ^ω of infinite words over a countable alphabet Σ which is called the prefix metric and defined as follows. For $u, v \in \Sigma^\omega$ and $u \neq v$ let $d(u, v) = 2^{-l_{\text{pref}}(u, v)}$ where $l_{\text{pref}}(u, v)$ is the first integer n such that the $(n+1)^{\text{th}}$ letter of u is different from the $(n+1)^{\text{th}}$ letter of v . The topology induced on Σ^ω by this metric is just the product topology of the discrete topology on Σ . For $s \in \Sigma^{<\omega}$, the set $N_s := \{\alpha \in \Sigma^\omega \mid s \prec \alpha\}$ is a basic clopen (i.e., closed and open) set of Σ^ω . More generally open sets of Σ^ω are in the form $W \cap \Sigma^\omega$, where $W \subseteq \Sigma^{<\omega}$.

When Σ is a finite alphabet, the prefix metric induces on Σ^ω the usual Cantor topology and Σ^ω is compact.

The Baire space ω^ω is equipped with the product topology of the discrete topology on ω . It is homeomorphic to $P_\infty := \{\alpha \in 2^\omega \mid \forall i \in \omega \exists j \geq i \ \alpha(j) = 1\} \subseteq 2^\omega$, via the map defined on ω^ω by $H(\beta) := 0^{\beta(0)}10^{\beta(1)}1 \dots$.

We define now the **Borel Hierarchy** on a topological space X :

Definition 1. The classes $\Sigma_\xi^0(X)$ and $\Pi_\xi^0(X)$ of the Borel Hierarchy on the topological space X are defined as follows:

$\Sigma_1^0(X)$ is the class of open subsets of X .

$\Pi_1^0(X)$ is the class of closed subsets of X .

And for any countable ordinal $\xi \geq 2$:

$\Sigma_\xi^0(X)$ is the class of countable unions of subsets of X in $\cup_{\gamma < \xi} \Pi_\gamma^0$.

$\Pi_\xi^0(X)$ is the class of countable intersections of subsets of X in $\cup_{\gamma < \xi} \Sigma_\gamma^0$.

As usual the ambiguous class Δ_ξ^0 is the class $\Sigma_\xi^0 \cap \Pi_\xi^0$.

Suppose now that $X \subseteq Y$; then $\Sigma_\xi^0(X) = \{A \cap X \mid A \in \Sigma_\xi^0(Y)\}$, and similarly for Π_ξ^0 , see [Kec95, Section 22.A]. Notice that we have defined the Borel classes $\Sigma_\xi^0(X)$ and $\Pi_\xi^0(X)$ mentioning the space X . However when the context is clear we will sometimes omit X and denote $\Sigma_\xi^0(X)$ by Σ_ξ^0 and similarly for the dual class.

The class of **Borel sets** is $\Delta_1^1 := \cup_{\xi < \omega_1} \Sigma_\xi^0 = \cup_{\xi < \omega_1} \Pi_\xi^0$, where ω_1 is the first uncountable ordinal.

For a countable ordinal α , a subset of Σ^ω is a Borel set of rank α iff it is in $\Sigma_\alpha^0 \cup \Pi_\alpha^0$ but not in $\cup_{\gamma < \alpha} (\Sigma_\gamma^0 \cup \Pi_\gamma^0)$.

We now define completeness with regard to reduction by continuous functions. For a countable ordinal $\alpha \geq 1$, a set $F \subseteq \Sigma^\omega$ is said to be a Σ_α^0 (respectively, Π_α^0)-complete set iff for any set $E \subseteq Y^\omega$ (with Y a finite alphabet): $E \in \Sigma_\alpha^0$ (respectively, $E \in \Pi_\alpha^0$) iff there exists a continuous function $f : Y^\omega \rightarrow \Sigma^\omega$ such that $E = f^{-1}(F)$. Σ_n^0 (respectively, Π_n^0)-complete sets, with n an integer ≥ 1 , are thoroughly characterized in [Sta86].

Recall that a set $X \subseteq \Sigma^\omega$ is a Σ_α^0 (respectively Π_α^0)-complete subset of Σ^ω iff it is in Σ_α^0 but not in Π_α^0 (respectively in Π_α^0 but not in Σ_α^0), [Kec95].

For example, the singletons of 2^ω are Π_1^0 -complete subsets of 2^ω . The set P_∞ is a well known example of a Π_2^0 -complete subset of 2^ω .

If Γ is a class of sets, then $\check{\Gamma} := \{\neg A \mid A \in \Gamma\}$ is the class of complements of sets in Γ . In particular, for every non-null countable ordinal α , $\check{\Sigma}_\alpha^0 = \Pi_\alpha^0$ and $\check{\Pi}_\alpha^0 = \Sigma_\alpha^0$.

We now introduce the Wadge hierarchy, which is a great refinement of the Borel hierarchy defined via reductions by continuous functions, [Wad83,Dup01].

Definition 2 (Wadge [Wad83]). Let X, Y be two finite alphabets. For $L \subseteq X^\omega$ and $L' \subseteq Y^\omega$, L is said to be Wadge reducible to L' ($L \leq_W L'$) iff there exists a continuous function $f : X^\omega \rightarrow Y^\omega$, such that $L = f^{-1}(L')$.

L and L' are Wadge equivalent iff $L \leq_W L'$ and $L' \leq_W L$. This will be denoted by $L \equiv_W L'$. And we shall say that $L <_W L'$ iff $L \leq_W L'$ but not $L' \leq_W L$.

A set $L \subseteq X^\omega$ is said to be self dual iff $L \equiv_W L^-$, and otherwise it is said to be non self dual.

The relation \leq_W is reflexive and transitive, and \equiv_W is an equivalence relation.

The equivalence classes of \equiv_W are called Wadge degrees.

The Wadge hierarchy WH is the class of Borel subsets of a set X^ω , where X is a finite set, equipped with \leq_W and with \equiv_W .

For $L \subseteq X^\omega$ and $L' \subseteq Y^\omega$, if $L \leq_W L'$ and $L = f^{-1}(L')$ where f is a continuous function from X^ω into Y^ω , then f is called a continuous reduction of L to L' . Intuitively it means that L is less complicated than L' because to check whether $x \in L$ it suffices to check whether $f(x) \in L'$ where f is a continuous function. Hence the Wadge degree of an ω -language is a measure of its topological complexity.

Notice that in the above definition, we consider that a subset $L \subseteq X^\omega$ is given together with the alphabet X .

We can now define the Wadge class of a set L :

Definition 3. Let L be a subset of X^ω . The Wadge class of L is :

$$[L] = \{L' \mid L' \subseteq Y^\omega \text{ for a finite alphabet } Y \text{ and } L' \leq_W L\}.$$

Recall that each Borel class Σ_α^0 and Π_α^0 is a Wadge class. A set $L \subseteq X^\omega$ is a Σ_α^0 (respectively Π_α^0)-complete set iff for any set $L' \subseteq Y^\omega$, L' is in Σ_α^0 (respectively Π_α^0) iff $L' \leq_W L$.

Theorem 4 (Wadge). Up to the complement and \equiv_W , the class of Borel subsets of X^ω , for a finite alphabet X , is a well ordered hierarchy. There is an ordinal $|WH|$, called the length of the hierarchy, and a map d_W^0 from WH onto $|WH| - \{0\}$, such that for all $L, L' \subseteq X^\omega$:

$$\begin{aligned} d_W^0 L < d_W^0 L' &\leftrightarrow L <_W L' \text{ and} \\ d_W^0 L = d_W^0 L' &\leftrightarrow [L \equiv_W L' \text{ or } L \equiv_W L'^-]. \end{aligned}$$

The Wadge hierarchy of Borel sets of **finite rank** has length ${}^1\varepsilon_0$ where ${}^1\varepsilon_0$ is the limit of the ordinals α_n defined by $\alpha_1 = \omega_1$ and $\alpha_{n+1} = \omega_1^{\alpha_n}$ for n a non negative integer, ω_1 being the first non countable ordinal. Then ${}^1\varepsilon_0$ is the first fixed point of the

ordinal exponentiation of base ω_1 . The length of the Wadge hierarchy of Borel sets in $\Delta_\omega^0 = \Sigma_\omega^0 \cap \Pi_\omega^0$ is the ω_1^{th} fixed point of the ordinal exponentiation of base ω_1 , which is a much larger ordinal. The length of the whole Wadge hierarchy of Borel sets is a huge ordinal, with regard to the ω_1^{th} fixed point of the ordinal exponentiation of base ω_1 . It is described in [Wad83,Dup01] by the use of the Veblen functions.

There are some subsets of the topological space Σ^ω which are not Borel sets. In particular, there exists another hierarchy beyond the Borel hierarchy, called the projective hierarchy. The first class of the projective hierarchy is the class Σ_1^1 of **analytic** sets. A set $A \subseteq \Sigma^\omega$ is analytic iff there exists a Borel set $B \subseteq (\Sigma \times Y)^\omega$, with Y a finite alphabet, such that $x \in A \leftrightarrow \exists y \in Y^\omega$ such that $(x, y) \in B$, where $(x, y) \in (\Sigma \times Y)^\omega$ is defined by: $(x, y)(i) = (x(i), y(i))$ for all integers $i \geq 0$.

A subset of Σ^ω is analytic if it is empty, or the image of the Baire space by a continuous map. The class of analytic sets contains the class of Borel sets in any of the spaces Σ^ω . Notice that $\Delta_1^1 = \Sigma_1^1 \cap \Pi_1^1$, where Π_1^1 is the class of co-analytic sets, i.e. of complements of analytic sets.

The ω -power of a finitary language V is always an analytic set because if V is finite and has n elements then V^ω is the continuous image of a compact set $\{0, 1, \dots, n-1\}^\omega$ and if V is infinite then there is a bijection between V and ω and V^ω is the continuous image of the Baire space ω^ω , [Sim92].

3 Topological complexity of ω -powers

We now state our first main result, showing that ω -powers exhibit a very surprising topological complexity.

Theorem 5 ([FL07]). *Let ξ be a non-null countable ordinal.*

- (a) *There is $A \subseteq 2^{<\omega}$ such that A^ω is Σ_ξ^0 -complete.*
- (b) *There is $A \subseteq 2^{<\omega}$ such that A^ω is Π_ξ^0 -complete.*

To prove Theorem 5, we use in [FL07] a level by level version of a theorem of Lusin and Souslin stating that every Borel set $B \subseteq 2^\omega$ is the image of a closed subset of the Baire space ω^ω by a continuous bijection, see [Kec95, p.83]. It is the following theorem, proved by Kuratowski in [Kur66, Corollary 33.II.1]:

Theorem 6. *Let ξ be a non-null countable ordinal, and $B \in \Pi_{\xi+1}^0(2^\omega)$. Then there is $C \in \Pi_1^0(\omega^\omega)$ and a continuous bijection $f : C \rightarrow B$ such that f^{-1} is Σ_ξ^0 -measurable (i.e., $f[U]$ is $\Sigma_\xi^0(B)$ for each open subset U of C).*

The existence of the continuous bijection $f : C \rightarrow B$ given by this theorem (without the fact that f^{-1} is Σ_ξ^0 -measurable) has been used by Arnold in [Arn83] to prove that every Borel subset of Σ^ω , for a finite alphabet Σ , is accepted by a non-ambiguous finitely branching transition system with Büchi acceptance condition. Notice that the sets of states of these transition systems are countable.

Our first idea was to code the behaviour of such a transition system. In fact this can be done on a part of ω -words of a special compact set $K_{0,0}$. However we have also to consider more general sets $K_{N,j}$ and then we need the hypothesis of the Σ_ξ^0 -measurability of the function f . The complete proof can be found in [FL07,FL08].

Notice that for the class Σ_2^0 , we need another proof, which uses a new operation which is very close to the erasing operation defined by Duparc in his study of the Wadge hierarchy, [Dup01]. We get the following result.

Theorem 7. *There is a context-free language $A \subseteq 2^{<\omega}$ such that $A^\omega \in \Sigma_2^0 \setminus \Pi_2^0$.*

Notice that it is easy to see that the set $2^\omega \setminus P_\infty$, which is the classical example of Σ_2^0 -complete set, is not an ω -power. The question is still open to know whether there exists a regular language L such that L^ω is Σ_2^0 -complete.

Recall that, for each non-null countable ordinal ξ , the class of Σ_ξ^0 -complete (respectively, Π_ξ^0 -complete) subsets of 2^ω forms a single *non self-dual* Wadge degree. Thus Theorem 5 provides also some Wadge degrees of ω -powers. More generally, it is natural to ask for the Wadge hierarchy of ω -powers. In the long version [FL08] of the conference paper [FL07] we get many more Wadge degrees of ω -powers.

In order to state these new results, we now recall the notion of difference hierarchy. (Recall that a countable ordinal γ is said to be even iff it can be written in the form $\gamma = \alpha + n$, where α is a limit ordinal and n is an even positive integer; otherwise the ordinal γ is said to be odd; notice that all limit ordinals are even ordinals.)

If $\eta < \omega_1$ and $(A_\theta)_{\theta < \eta}$ is an increasing sequence of subsets of some space X , then we set

$$D_\eta[(A_\theta)_{\theta < \eta}] := \{x \in X \mid \exists \theta < \eta \ x \in A_\theta \setminus \bigcup_{\theta' < \theta} A_{\theta'} \text{ and the parity of } \theta \text{ is opposite to that of } \eta\}.$$

If moreover $1 \leq \xi < \omega_1$, then we set :

$$D_\eta(\Sigma_\xi^0) := \{D_\eta[(A_\theta)_{\theta < \eta}] \mid \text{for each } \theta < \eta \ A_\theta \text{ is in the class } \Sigma_\xi^0\}.$$

Recall that for each non null countable ordinal ξ , the sequence $(D_\eta(\Sigma_\xi^0))_{\eta < \omega_1}$ is strictly increasing for the inclusion relation and that for each $\eta < \omega_1$ it holds that $D_\eta(\Sigma_\xi^0) \subseteq \Delta_{\xi+1}^0$. Moreover for each $\eta < \omega_1$ the class $D_\eta(\Sigma_\xi^0)$ is a Wadge class and the class of $D_\eta(\Sigma_\xi^0)$ -complete subsets of 2^ω forms a single *non self-dual* Wadge degree.

Theorem 8.

1. *Let $1 \leq \xi < \omega_1$. Then there is $A \subseteq 2^{<\omega}$ such that A^ω is $\check{D}_2(\Sigma_\xi^0)$ -complete.*
2. *Let $3 \leq \xi < \omega_1$ and $1 \leq \theta < \omega_1$. Then there is $A \subseteq 2^{<\omega}$ such that A^ω is $\check{D}_{\omega^\theta}(\Sigma_\xi^0)$ -complete.*

Notice that for each ordinal ξ such that $3 \leq \xi < \omega_1$ we get uncountably many Wadge degrees of ω -powers of the same Borel rank $\xi + 1$. This confirms the great complexity of these ω -languages.

However the problem is still open to determine completely the Wadge hierarchy of ω -powers.

We now come to the effectiveness questions. It is natural to wonder whether the ω -powers obtained above are effective. For instance could they be obtained as ω -powers of recursive languages ?

In the paper [FL08] we prove effective versions of the results presented above. Using tools of effective descriptive set theory, such Kleene recursion Theorem and the notion of Borel codes, we first prove an effective version of Kuratowski's Theorem 6. Then we use it to prove the following effective version of Theorem 5, where Σ_ξ^0 and Π_ξ^0 denote classes of the hyperarithmetical hierarchy and ω_1^{CK} is the first non-recursive ordinal, usually called the Church-Kleene ordinal.

Theorem 9. *Let ξ be a non-null ordinal smaller than ω_1^{CK} .*

(a) *There is a recursive language $A \subseteq 2^{<\omega}$ such that $A^\omega \in \Sigma_\xi^0 \setminus \Pi_\xi^0$.*

(b) *There is a recursive language $A \subseteq 2^{<\omega}$ such that $A^\omega \in \Pi_\xi^0 \setminus \Sigma_\xi^0$.*

Remark 10. *If $A \subseteq 2^{<\omega}$ is a recursive language, then the ω -power A^ω is an effective analytic set, i.e. a (lightface) Σ_1^1 -set. And the supremum of the set of Borel ranks of Borel effective analytic sets is the ordinal γ_2^1 . This ordinal is defined by Kechris, Marker, and Sami in [KMS89] and it is proved to be strictly greater than the ordinal δ_2^1 which is the first non Δ_2^1 ordinal. Thus the ordinal γ_2^1 is also strictly greater than the first non-recursive ordinal ω_1^{CK} . Thus Theorem 9 does not give the complete answer about the Borel hierarchy of ω -powers of recursive languages. Indeed there could exist some ω -powers of recursive languages of Borel ranks greater than ω_1^{CK} , but of course smaller than the ordinal γ_2^1 .*

4 Concluding remarks

The question naturally arises to know what are all the possible infinite Borel ranks of ω -powers of finitary languages belonging to some natural class like the class of context free languages (respectively, languages accepted by stack automata, recursive languages, recursively enumerable languages, ...).

We know from [Fin06] that there are ω -languages accepted by Büchi 1-counter automata of every Borel rank (and even of every Wadge degree) of an effective analytic set. Every ω -language accepted by a Büchi 1-counter automaton can be written as a finite union $L = \bigcup_{1 \leq i \leq n} U_i \hat{\ } V_i^\omega$, where for each integer i , U_i and V_i are finitary languages accepted by 1-counter automata. And the supremum of the set of Borel ranks of effective analytic sets is the ordinal γ_2^1 . From these results it seems plausible that there exist some ω -powers of languages accepted by 1-counter automata which have Borel

ranks up to the ordinal γ_2^1 , although these languages are located at the very low level in the complexity hierarchy of finitary languages.

Another interesting question would be to determine completely the Wadge hierarchy of ω -powers. A simpler open question is to determine the Wadge hierarchy of ω -powers of regular languages. The second author has given in [Lec05] a few Wadge degrees of ω -powers of regular languages. Notice however that even the question to determine the Wadge degrees of ω -powers of regular languages in the class Δ_2^0 is still open.

References

- Arn83. A. Arnold, Topological Characterizations of Infinite Behaviours of Transition Systems, Automata, Languages and Programming, J. Diaz Ed., Lecture Notes in Computer Science, Volume 154, Springer, 1983, p. 28-38.
- ABB96. J-M. Autebert, J. Berstel and L. Boasson, Context Free Languages and Pushdown Automata, in Handbook of Formal Languages, Vol 1, Springer Verlag 1996.
- Dup01. J. Duparc, Wadge Hierarchy and Veblen Hierarchy: Part 1: Borel Sets of Finite Rank, Journal of Symbolic Logic, Vol. 66, no. 1, 2001, p. 56-86.
- DF07. J. Duparc and O. Finkel, An ω -Power of a Context-Free Language Which Is Borel Above Δ_ω^0 , in the Proceedings of the International Conference Foundations of the Formal Sciences V : Infinite Games, November 26th to 29th, 2004, Bonn, Germany, Volume 11 of Studies in Logic, College Publications at King's College, London, 2007, p. 109-122.
- Fin01. O. Finkel, Topological Properties of Omega Context Free Languages, Theoretical Computer Science, Vol. 262 (1-2), July 2001, p. 669-697.
- Fin03. O. Finkel, Borel Hierarchy and Omega Context Free Languages, Theoretical Computer Science, Vol. 290 (3), 2003, p. 1385-1405.
- Fin04. O. Finkel, An omega-Power of a Finitary Language Which is a Borel Set of Infinite Rank, Fundamenta Informaticae, Volume 62 (3-4), 2004, p. 333-342.
- Fin06. O. Finkel, Borel Ranks and Wadge Degrees of Omega Context Free Languages, Mathematical Structures in Computer Science, Volume 16 (5), 2006, p. 813-840.
- FL07. O. Finkel and D. Lecomte, There Exist some ω -Powers of Any Borel Rank. In the Proceedings of the 16th EACSL Annual International Conference on Computer Science and Logic, CSL 2007, Lausanne, Switzerland, September 11-15, 2007, Lecture Notes in Computer Science, Volume 4646, Springer, 2007, p. 115-129.
- FL08. O. Finkel and D. Lecomte, Classical and Effective Descriptive Complexities of omega-Powers, preprint, 2008, available from <http://fr.arxiv.org/abs/0708.4176>.
- HU69. J.E. Hopcroft and J.D. Ullman, Formal Languages and their Relation to Automata, Addison-Wesley Publishing Company, Reading, Massachusetts, 1969.
- Kec95. A.S. Kechris, Classical Descriptive Set Theory, Springer-Verlag, 1995.
- KMS89. A. S. Kechris, D. Marker, and R. L. Sami, Π_1^1 Borel Sets, The Journal of Symbolic Logic, Volume 54 (3), 1989, p. 915-920.
- Kur66. K. Kuratowski, Topology, Vol. 1, Academic Press, New York 1966.
- Lec01. D. Lecomte, Sur les Ensembles de Phrases Infinies Constructibles a Partir d'un Dictionnaire sur un Alphabet Fini, Séminaire d'Initiation a l'Analyse, Volume 1, année 2001-2002.
- Lec05. D. Lecomte, Omega-Powers and Descriptive Set Theory, Journal of Symbolic Logic, Volume 70 (4), 2005, p. 1210-1232.
- LT94. H. Lescow and W. Thomas, Logical Specifications of Infinite Computations, In: "A Decade of Concurrency" (J. W. de Bakker et al., eds), Springer LNCS 803 (1994), 583-621.

- Mos80. Y. N. Moschovakis, Descriptive Set Theory, North-Holland, Amsterdam 1980.
- Niw90. D. Niwinski, Problem on ω -Powers posed in the Proceedings of the 1990 Workshop "Logics and Recognizable Sets" (Univ. Kiel).
- PP04. D. Perrin and J.-E. Pin, Infinite Words, Automata, Semigroups, Logic and Games, Volume 141 of Pure and Applied Mathematics, Elsevier, 2004.
- Sim92. P. Simonnet, Automates et Théorie Descriptive, Ph. D. Thesis, Université Paris 7, March 1992.
- Sta86. L. Staiger, Hierarchies of Recursive ω -Languages, Jour. Inform. Process. Cybernetics EIK 22 (1986) 5/6, 219-241.
- Sta97a. L. Staiger, ω -Languages, Chapter of the Handbook of Formal Languages, Vol 3, edited by G. Rozenberg and A. Salomaa, Springer-Verlag, Berlin, 1997.
- Sta97b. L. Staiger, On ω -Power Languages, in New Trends in Formal Languages, Control, Coperation, and Combinatorics, Lecture Notes in Computer Science 1218, Springer-Verlag, Berlin 1997, 377-393.
- Tho90. W. Thomas, Automata on Infinite Objects, in: J. Van Leeuwen, ed., Handbook of Theoretical Computer Science, Vol. B (Elsevier, Amsterdam, 1990), p. 133-191.
- Wad83. W. W. Wadge. Reducibility and Determinateness in the Baire Space, PhD thesis, University of California, Berkeley, 1983.