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A New Range-Reduction Algorithm

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Abstract—Range-reduction is a key point for getting accurate elementary function routines. We introduce a new algorithm that is fast for input arguments belonging to the most common domains, yet accurate over the full double-precision range.

Index Terms—Range-reduction, elementary function evaluation, floating-point arithmetic.

1 INTRODUCTION

Algorithms for the evaluation of elementary functions give correct results only if the argument is within a given small interval, usually centered at zero. To evaluate an elementary function \( f(x) \) for any \( x \), it is necessary to find some “transformation” that makes it possible to deduce \( f(x) \) from some value \( g(x^*) \), where:

- \( x^* \), called the reduced argument, is deduced from \( x \);
- \( x^* \) belongs to the convergence domain of the algorithm implemented for the evaluation of \( g \).

In practice, range-reduction needs care for the trigonometric functions. With these functions, \( x^* \) is equal to \( x - kC \), where

\[ k \text{ is an integer and } C \text{ an integer multiple of } \pi/4. \]

Also of potential interest is the case \( C = \ln(2) \) for the implementation of the exponential function.

A poor range-reduction method may lead to catastrophic accuracy problems when the input argument is large or close to an integer multiple of \( C \). It is easy to understand why a poor range-reduction algorithm gives inaccurate results. The naïve method consists of performing the computations

\[ k = \left\lfloor \frac{x}{C} \right\rfloor \]
\[ x^* = x - kC, \]

using machine precision. When \( kC \) is close to \( x \), almost all the accuracy, if not all, is lost when performing the subtraction \( x - kC \). For instance, if \( C = \pi/2 \) and \( x = 8248.251512 \), the correct value of \( x^* \) is \(-2.14758367 \times 10^{12}\), and the corresponding value of \( k \) is 5.251. Directly computing \( x - k\pi/2 \) on a calculator with 10-digit decimal arithmetic (assuming rounding to the nearest and replacing \( \pi/2 \) by the nearest exactly representable number), then one gets \(-1.0 \times 10^{-6}\). Hence, such a poor range-reduction would lead to a computed value of \( \cos(x^*) \) equal to \(-1.0 \times 10^{-6}\), whereas the correct value is \(-2.14758367 \times 10^{12}\).

A first solution to overcome the problem consists of using arbitrary-precision arithmetic, but this may make the computation much slower. Moreover, it is not that easy to predict on the fly the precision with which the computation should be performed.

Most common input arguments to the trigonometric functions are small (say, less than 8) or sometimes medium (say, between 8 and approximately 2^{64}). They are rarely huge (say, greater than 2^{64}). We want to design methods that are fast for the frequent cases, and accurate for all cases. A rough estimate, based on SUN fdlibm library, is that the cost of trigonometric range-reduction—when reduction is necessary—is approximately one third of the total function evaluation cost.

First, we describe Payne and Hanek’s method [11], which provides an accurate range-reduction, but has the drawback of being fairly expensive in term of operations; this method is very commonly implemented; it is used in the SUN fdlibm library in particular.

To know with which precision the intermediate calculations must be carried on to get an accurate result, one must know the worst cases, that is, the input arguments that are hardest to reduce. Also, to estimate the average performance of the algorithms (and to tune them so that these performances are good), one must have at least a rough estimate of the statistical distribution of the reduced arguments. These two problems are dealt with at the end of this section.

In the second section, we present our algorithm dedicated to the reduction of small and medium size arguments. In the third section, we compare our method with some other available methods, which justifies the use of our algorithm for small and medium size arguments.

1.1 The Payne and Hanek Reduction Algorithm

We assume in this section that we want to perform range-reduction for the trigonometric functions, with \( C = \pi/4 \), and that the convergence domain of the algorithm used for...
evaluating the functions contains\(^1\) \( I = [0, \pi/4] \). An adaptation to other cases is straightforward.

From an input argument \( x \), we want to find the reduced argument \( x^* \) and an integer \( k \) that satisfy:

\[
k = \left\lfloor \frac{4}{\pi} x \right\rfloor \quad x^* = \pi \left( \frac{4}{\pi} x - k \right).
\]

Once \( x^* \) is known, it suffices to know \( k \mod 8 \) to calculate \( \sin(x) \) or \( \cos(x) \) from \( x^* \). If \( x \) is large or if \( x \) is very close to a multiple of \( \pi/4 \), the direct use of (1) to determine \( x^* \) may require the knowledge of \( 4/\pi \) with very large precision and a cost-expensive multiple-precision compilation if we wish the range-reduction to be accurate.

Now, let us present Payne and Hanek's reduction method [11], [12]. Assume an \( n \)-bit mantissa, radix 2 floating-point format (the number of bits \( n \) includes the possible hidden bit; for instance, with an IEEE double-precision format, \( n = 53 \)). Let \( x \) be the positive floating-point argument to be reduced and let \( e \) be its unbiased exponent, so

\[
x = X \times 2^{-n+1},
\]

where \( X \) is an \( n \)-bit integer satisfying \( 2^{n-1} \leq X < 2^n \).

We can assume \( e \geq -1 \) (since, if \( e < -1 \), no reduction is necessary). Let

\[
\alpha_0\alpha_{-1}\alpha_{-2}\alpha_{-3}\alpha_{-4}\alpha_{-5} \ldots
\]

be the infinite binary expansion of \( \alpha = 4/\pi \) and define an integer parameter \( p \) used to specify the required accuracy of the range-reduction. Then, rewrite \( \alpha = 4/\pi \) as

\[
\text{Left}(e, p) \times 2^{2-e+p} + (\text{Middle}(e, p) + \text{Right}(e, p)) \times 2^{-n-e-1-p},
\]

where

\[
\begin{align*}
\text{Left}(e, p) &= 0 \quad \text{if } e < n + 2 \\
\text{Middle}(e, p) &= \alpha_0\alpha_{-1}\ldots\alpha_{n-e+2} \text{ otherwise,} \\
\text{Right}(e, p) &= 0,\alpha_{n-e-2}\ldots\alpha_{n-e-3-p} \ldots
\end{align*}
\]

Fig. 1 shows the splitting of the binary expansion of \( \alpha \).

The basic idea of the Payne-Hanek reduction method is to notice that, if \( p \) is large enough, \( \text{Middle}(e, p) \) contains the only bits of \( \alpha = 4/\pi \) that matter for the range-reduction. Since

\[
\frac{4}{\pi} x = \text{Left}(e, p) \times X \times 8 \\
+ \text{Middle}(e, p) \times X \times 2^{-2n-p} \\
+ \text{Right}(e, p) \times X \times 2^{-2n-p},
\]

1. In practice, we can reduce to an interval of size slightly larger than \( C \) to facilitate the reduction.
1.3 Statistical Distribution of the Reduced Arguments

Now, let us turn to the statistical distribution of reduced arguments.

We assume that $C$ is a positive fractional multiple of $\pi$ or $\ln(2)$. Let $e_{\min}$ and $e_{\max}$ be two rational integers such that $2^{e_{\min}} \leq C/2 < 2^{e_{\min} + 1}$ and $e_{\min} \leq e_{\max}$.

Let $p \in \mathbb{N}$ such that $2^{-p+1} \leq C$, our aim is to estimate the number of floating-point numbers $x$ with $n$-bit mantissas and exponents between $e_{\min}$ and $e_{\max}$ such that

$$|x \mod^* C| < 2^{-p},$$

(2)

where $x \mod^* C$ is defined as the unique number $y \in [-C/2, +C/2)$ such that $y = x - kC$, where $k$ is an integer.

Let $E$ be a rational integer such that $e_{\min} \leq E \leq e_{\max}$. As $2^{-p+1} \leq C$, we have $2^{-p} < 2^{e_{\min} + 1} \leq 2^{E+1}$. Therefore, $2^{-p} < 2^E$, i.e., $p + E \geq 0$.

We start by estimating the number of floating-point numbers $x$ with $n$-bit mantissas and exponent $E$ that satisfy (2). Hence, we search for the $j \in \mathbb{N}$, $2^{n-1} \leq j \leq 2^n - 1$, such that the inequality

$$|kC - \frac{j}{2^n - 1}2^E| < 2^{-p}$$

(3)

has solutions in $k \in \mathbb{Z}$. Such $k$ necessarily satisfy

$$\frac{1}{C} \left( \frac{1}{2^p} + \frac{j}{2^n - 1}2^E \right) < k < \frac{1}{C} \left( \frac{1}{2^p} + \frac{j}{2^n - 1}2^E \right).$$

(4)

We note that, as $p + E \geq 0$ and $j \geq 2^{n-1}$, the left-hand side of (4) is positive. Hence,

$$\max \left\{ \frac{1}{C} \left( \frac{1}{2^p} + \frac{j}{2^n - 1}2^E \right) \right\} \leq k \leq \frac{1}{C} \left( \frac{1}{2^p} + \frac{2^{E+1} - 2^E}{2^n - 1} \right),$$

(5)

since $2^{n-1} \leq j \leq 2^n - 1$ and these inequalities are sharp since the upper bound in (4) is irrational and the lower bound is either zero or an irrational number. The number of possible $k$ is exactly

$$N_E = M_E - m_E + 1.$$  

(6)

Inequality (3) is equivalent to

$$|kC2^{n-1} - j| < 2^{n-1-p}E.$$  

(7)

Hence, for every $k$ satisfying (5), there are exactly

$$\min(2^{n-1} - |kC2^{n-1} - j| + 1, 2^{n-1} - |kC2^{n-1} - j| - 2^{n-1-p}E)$$

(8)

integers $j$ solutions since the numbers $kC2^{n-1} - 2^{n-1-p}E$ and $kC2^{n-1} - 2^{n-1-p}E$ are irrational (we saw before that $k \neq 0$).

As $2^{-p+1} \leq C$, if $k \geq m_E + 1$, we have

$$2^{n-1} \leq |kC2^{n-1} - 2^{n-1-p}E|$$

and, if $k \leq M_E - 1$, we have

$$2^{n-1} - 2^{n-1-p}E \geq |kC2^{n-1} - 2^{n-1-p}E|.$$  

Now, to analyze (8), we have to distinguish two cases.

**First case:** $2^{n-1-p}E \geq 1/2$, i.e., $n - E \geq p$.

This case is the easy one and (7) yields the conclusion. For every $k$, $m_E + 1 \leq k \leq M_E - 1$, there are exactly $2^{n-1-p}E$ integer solutions $j$ since the numbers $kC2^{n-1} - 2^{n-1-p}E$ and $kC2^{n-1} - 2^{n-1-p}E$ are irrational. When $k \in \{m_E, M_E\}$, we can only say that there are at least 1 and at most $2^{n-1-p}E$ integer solutions $j$. Notice that these solutions can easily be enumerated by a program. Therefore, the number of floating-point numbers $x$ with $n$-bit mantissas and exponent $E$ that satisfy (2) is upper bounded by $N_E2^{n-p}E$, and lower bounded by $(N_E - 2)2^{n-p}E + 2$.

**Second case:** $2^{n-1-p}E < 1/2$, i.e., $n - E < p$.

We need results about uniform distribution of sequences [8] that we briefly recall now.

For a real number $x$, $\lfloor x \rfloor$ denotes the fractional part of $x$, i.e., $\{x\} = x - \lfloor x \rfloor$ and $\|x\|$ denotes the distance from $x$ to the nearest integer, namely,

$$\|x\| = \min_{n \in \mathbb{Z}} |x - n| = \min(\{x\}, 1 - \{x\}).$$

Let us recall the following definitions from [8].

**Definition 1.** Let $(x_n)_{n \geq 1}$ be a given sequence of real numbers. Let $N$ be a positive integer.

For a subset $E$ of $[0, 1)$, the counting function $A(E; N; (x_n))$ is the number of terms $x_n$, $1 \leq n \leq N$, for which $\{x_n\} \in E$.

Let $y_1, \ldots, y_N$ be a finite sequence of real numbers. The number

$$D_N((y_n)) = \sup_{0 \leq a < b \leq 1} \frac{A((a, b); N; (y_n)) - (b - a)}{N}$$

is called the discrepancy of the sequence $y_1, \ldots, y_N$. For an infinite sequence $(x_n)$ of real numbers (or for a finite sequence containing at least $N$ terms), $D_N((x_n))$ is meant to be the discrepancy of the initial segment formed by the first $N$ terms of $(x_n)$.

Thus, in particular, the number of values $x_n$ with $1 \leq n \leq N$ satisfying $\{x_n\} \in [a, b]$, for any $0 \leq a < b \leq 1$, is bounded from above by $N((\{a\} + D_N((x_n))))$. Hence, the number of values $kC2^{n-1} - E$, with $m_E \leq k \leq M_E$, that satisfy (7), i.e., that satisfy $0 \leq \{kC2^{n-1} - 2^{n-1-p}E\} < 2^{n-1-p}E$ or $1 - 2^{n-1-p}E < \{kC2^{n-1} - 2^{n-1-p}E\} < 1$ is bounded from above by $N_E2^{n-p}E + 2D_N((kC2^{n-1}E)))$.

**Definition 2.** Let $\mu$ be a positive real number or infinity. The irrational number $\alpha$ is said to be of type $\mu$ if $\mu$ is the supremum of all $\gamma$ for which

$$\liminf_{n \to \infty} \frac{q_n \gamma \|q_n\|}{n^\mu} = 0.$$

Theorem 3.2 from [8, chapter 2] states the following result:

**Theorem 1.** Let $\alpha$ be of finite type $\mu$. Then, for every $\varepsilon > 0$, the discrepancy $D_N(u)$ of $u = (n\alpha)$ satisfies

$$D_N(u) = O(N^{1/(\mu + \varepsilon)}).$$
Proposition 1. Let $C$ be a positive fractional multiple of $\pi$ or $\ln(2)$. Let $\epsilon_\text{min}$ and $\epsilon_\text{max}$ be two rational integers such that $2^{\epsilon_\text{min}} \leq C/2 < 2^{\epsilon_\text{max}+1}$ and $\epsilon_\text{min} \leq \epsilon_\text{max}$. Let $p \in \mathbb{N}$ such that $2^{-p} \leq C/2$. The number $\nu_E$ of floating-point numbers $x$ with $n$-bit mantissas and exponent $E$ that satisfy (2) is upper bounded by $2^{n-p-E}(N_E + O(N_E(e^{\delta/2})))$ for every $\epsilon > 0$.

From this theorem, we can deduce the following result.

### Proposition 1

Let $C$ be a positive fractional multiple of $\pi$ or $\ln(2)$. Let $\epsilon_\text{min}$ and $\epsilon_\text{max}$ be two rational integers such that $2^{\epsilon_\text{min}} \leq C/2 < 2^{\epsilon_\text{max}+1}$ and $\epsilon_\text{min} \leq \epsilon_\text{max}$. Let $p \in \mathbb{N}$ such that $2^{-p} \leq C/2$. The number $\nu_E$ of floating-point numbers $x$ with $n$-bit mantissas and exponent $E$ that satisfy (2) is upper bounded by $2^{n-p-E}(N_E + O(N_E(e^{\delta/2})))$ for every $\epsilon > 0$.

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2 A NEW HIGH-RADIX REDUCTION METHOD

In this section, we assume that we perform range-reduction for the trigonometric functions, with \( C = \frac{\pi}{2} \). Extension to other values of \( C \) (such as a fractional multiple of \( \pi \)—still for the trigonometric functions—or a fractional multiple of \( \ln(2) \)—for the exponential function) is straightforward.

As stated before, our general philosophy is that we must give results that are:

1. always correct, even for rare cases;
2. computed as quickly as possible for frequent cases.

A way to deal with these requirements is to build a fast algorithm for input arguments with a small exponent and to use a slower yet still accurate algorithm for input argument with a large exponent.

2.1 Medium-Size Arguments (in \([8, 2^{63} - 1]\))

To do so, in the following, we focus on input arguments with a “reasonably small” exponent. More precisely, we assume that the double-precision input argument \( x \) has absolute value less than \( 2^{63} - 1 \). For larger arguments, we assume that Payne and Hanek’s method will be used or that \( x \mod 2^C \) will be computed using multiple-precision arithmetic. For straightforward symmetry reasons, we can assume that \( x \) is positive. We also assume that \( x \) is greater than or equal to 8. We then proceed as follows:

1. We define \( I(x) \) as \( x \) rounded to the nearest integer. Then, \( x \) is split into its residual part \( \rho(x) = x - I(x) \) and \( I(x) \), which is split into eight 7-bit parts \( I_i(x) \) for \( 0 \leq i \leq 7 \) as follows:

\[
\begin{align*}
I_7(x) &= I(2^{56} x), \\
I_6(x) &= I(2^{48} (x - (2^{56} I_7(x))) ), \\
I_5(x) &= I(2^{40} (x - (2^{56} I_7(x)) + 2^{48} I_6(x))) , \\
I_4(x) &= I(2^{32} (x - \sum_{i=3}^{7} 2^{8 i} I_i(x)) ), \\
I_3(x) &= I(2^{24} (x - \sum_{i=1}^{7} 2^{8 i} I_i(x)) ), \\
I_2(x) &= I(2^{16} (x - \sum_{i=1}^{7} 2^{8 i} I_i(x)) ), \\
I_1(x) &= I(2^{8} (x - \sum_{i=1}^{7} 2^{8 i} I_i(x)) ), \\
I_0(x) &= I(x - \sum_{i=0}^{7} 2^{8 i} I_i(x)), \\
\rho(x) &= x - \sum_{i=0}^{7} 2^{8 i} I_i(x),
\end{align*}
\]

so that

\[ x = 2^{56} I_7(x) + 2^{48} I_6(x) + \ldots + 2^{8} I_1(x) + I_0(x) + \rho(x). \]

Note that \( \rho(x) \) is exactly representable in double-precision and that, for \( x \geq 2^{52} \), we have \( \rho(x) = 0 \) and \( I(x) = x \). Also, since \( x \geq 8 \), the last mantissa bit of \( \rho(x) \) has a weight greater than or equal to \( 2^{-49} \).

**Important remark.** One could get a very similar algorithm, certainly easier to understand, by replacing the values \( I_k(x) \) by the values \( J_k(x) \) defined as

\[
\begin{align*}
J_0(x) &\quad \text{contains bits 0 to 7 of } I(x), \\
J_1(x) &\quad \text{contains bits 8 to 15 of } I(x), \\
J_2(x) &\quad \text{contains bits 16 to 23 of } I(x), \\
J_3(x) &\quad \text{contains bits 24 to 31 of } I(x), \\
J_4(x) &\quad \text{contains bits 32 to 39 of } I(x), \\
J_5(x) &\quad \text{contains bits 40 to 47 of } I(x), \\
J_6(x) &\quad \text{contains bits 48 to 55 of } I(x), \\
J_7(x) &\quad \text{contains bits 56 to 63 of } I(x),
\end{align*}
\]

but that would lead to tables twice as large as the ones required by our algorithm. Indeed, the values \( I_0 \) up to \( I_7 \) are stored on 8 bits each, but the sign bit will not be used and, thus, only 7 bits are necessary to index the tables.

The general idea behind our algorithm is to compute first

\[
S(x) = \left( (I_0(x)) \mod \frac{\pi}{2} + \left( 2^8 I_1(x) \right) \right) \mod \frac{\pi}{2} + \left( 2^{16} I_3(x) \right) \mod \frac{\pi}{2} + \left( 2^{24} I_5(x) \right) \mod \frac{\pi}{2} + \rho(x).
\]

It holds that \( x - S(x) \) is a multiple of \( \frac{\pi}{2} \) and \( S(x) \) will be smaller than \( x \), but, in general, \( S(x) \) will not be the desired reduced argument: A second, simpler reduction step will be necessary. In practice, the various possible values of \( \left( \left( 2^{39} I_6(x) \right) \right) \mod \frac{\pi}{2} \) are stored in tables as a sum of two or three floating-point numbers.

As mentioned above, our goal is to always provide correct results even for the worst case for which we lose 61 bits of accuracy. Then, we need to store \( (I(x)) \mod \frac{\pi}{2} \) with at least

\[ 61 \text{ (leading zeros)} + 53 \text{ (nonzero significant bits)} + g \text{ (extra guard bits)} = 114 + g \text{ bits}. \]

To reach that precision (with a value of \( g \) equal to 39, which will be deduced in the following), all the numbers \( \left( \left( 2^{39} I_6(x) \right) \right) \mod \frac{\pi}{2} \), which belong to \([-1, 1]\), are stored in tables as the sum of three double-precision numbers:

\[
\begin{align*}
PT_{hi}(i, w) \quad &\text{is the multiple of } 2^{-49} \text{ that is}\quad \text{closest to } (2^w u) \mod \frac{\pi}{2} \text{ } \quad \text{in } \left( 2^{-49} \right) \\
T_{med}(i, w) \quad &\text{is the multiple of } 2^{-99} \text{ that is}\quad \text{closest to } (2^w u) \mod \frac{\pi}{2} \text{ } \quad \text{in } \left( 2^{-99} \right) \\
T_{lo}(i, w) \quad &\text{is the double-precision number } \quad \text{that is closest to}\quad (2^w u) \mod \frac{\pi}{2} - T_{hi}(i, w) - T_{med}(i, w),
\end{align*}
\]

where \( w \) is a 7-bit nonnegative integer.

Note that \( T_{hi}(i, w) = T_{med}(i, w) = T_{lo}(i, w) = 0 \) for \( w = 0 \). The three tables \( T_{hi}, T_{med}, \) and \( T_{lo} \) need 10 address bits. The total amount of memory required by
Define the rounding of precision, which corresponds to the absolute value is less than or equal to \( T_h(i, w) \). This table is useful for these values of \((i, w)\) allows us to get slightly sharper bounds, given in Table 1.

**TABLE 1**

<table>
<thead>
<tr>
<th>( \max_{w} T_h(i, w) )</th>
<th>( \max_{w} T_m(w) )</th>
<th>( \max_{w} T_o(i, w) )</th>
</tr>
</thead>
<tbody>
<tr>
<td>0.7846901 \ldots</td>
<td>0.397607 \ldots ( \times 2^{-50} )</td>
<td>0.998214 \ldots ( \times 2^{-199} )</td>
</tr>
</tbody>
</table>

the numbers are 3 \( \times 2^{10} \cdot 8 = 24 \) Kbytes. From the definitions, one can easily deduce \( |T_m(i, w)| \leq 2^{-50} \) and \( |T_o(i, w)| \leq 2^{-100} \). The sum \( T_h(i, w) + T_m(i, w) + T_o(i, w) \) approximates \((2^4 w) \mod \pi/2 \) with 153 bits of precision, which corresponds to \( g = 39 \). Computing \( T_h(i, w) \), \( T_m(i, w) \), and \( T_o(i, w) \) for the 1,024 different possible values of \((i, w)\) allows us to get slightly sharper bounds, given in Table 1.

2. Define

\[
S_h(x) = \left( \sum_{i=0}^{7} \text{sign}(I_i(x))T_h(i, |I_i(x)|) \right) + \rho(x).
\]

Its absolute value is bounded by \( 2\pi + \frac{1}{2} \), which is less than 8. Since \( S_h(x) \) is a multiple of \( 2^{-49} \) and has absolute value less than 8, it is exactly representable in double-precision floating-point arithmetic (it is even representable with 52 bits only). Therefore, with a correctly rounded arithmetic (such as the one provided on any system that complies with the IEEE-754 standard for floating-point arithmetic), it will be exactly computed, without any rounding error. Also, consider

\[
\begin{align*}
S_m(x) &= \sum_{i=0}^{7} \text{sign}(I_i(x))T_m(i, |I_i(x)|). \\
S_o(x) &= \sum_{i=0}^{7} \text{sign}(I_i(x))T_o(i, |I_i(x)|).
\end{align*}
\]

The number \( S_m(x) \) is a multiple of \( 2^{-99} \) and its absolute value is less than \( 2^{-47} \). Hence, it is exactly representable, and exactly computed, in double-precision floating-point arithmetic. \( |S_o| \) is less than \( 2^{-97} \) and, if \( S_o \) is computed with round-to-nearest arithmetic as a balanced binary tree of additions:

\[
\begin{align*}
&\left( \text{sign}(I_0(x))T_o(0, |I_0(x)|) \right) \\
+ &\text{sign}(I_1(x))T_o(1, |I_1(x)|) \\
+ &\text{sign}(I_2(x))T_o(2, |I_2(x)|) \\
+ &\text{sign}(I_3(x))T_o(3, |I_3(x)|) \\
+ &\left( \text{sign}(I_4(x))T_o(4, |I_4(x)|) \right) \\
+ &\text{sign}(I_5(x))T_o(5, |I_5(x)|) \\
+ &\text{sign}(I_6(x))T_o(6, |I_6(x)|) \\
+ &\text{sign}(I_7(x))T_o(7, |I_7(x)|)
\end{align*}
\]

then the rounding error is less than \( 3 \times 2^{-151} \). For each of the values \( T_o(i, I_i(x)) \), the fact that it rounded to the nearest yields an accumulated error (for these eight values) less than \( 8 \times 2^{-151} \). Thus, the absolute error on \( S_o(x) \) is less than or equal to \( 8 \times 2^{-154} + 3 \times 2^{-151} = 2^{-149} \).

Since \( S_h(x) + S_m(x) \) is exactly computed, the number \( S(x) = S_h(x) + S_m(x) + S_o(x) \) is equal to \( x \) minus an integer multiple of \( \pi/2 \) plus an error bounded by \( 2^{-149} \).

And yet, \( S(x) \) may not be the final reduced argument since its absolute value may be significantly larger than \( \pi/4 \). We therefore may have to add or subtract a multiple of \( \pi/2 \) from \( S(x) \) to get the final result and straightforward calculations show that this multiple can only be \( k\pi/2 \) with \( k = 1, 2, 3, \) or 4.

2.2 Small Arguments (Smaller than 8)

Define \( C_u(k) \), for \( k = 1, 2, 3, 4 \), as the multiple of \( 2^{-49} \) that is closest to \( k\pi/2 \). \( C_u(k) \) is exactly representable as a double-precision number. Define \( C_m(k) \) as the multiple of \( 2^{-99} \) that is closest to \( k\pi/2 - C_u(k) \) and \( C_o(k) \) as the double-precision number that is closest to \( k\pi/2 - C_u(k) - C_m(k) \).

We now proceed as follows:

- If \( |S_h(x)| \leq \pi/4 \), then we define
  \[
  R_h(x) = S_h(x), \\
  R_m(x) = S_m(x), \\
  R_o(x) = S_o(x).
  \]

- Else, let \( k_x \) be such that \( C_u(k_x) \) is closest to \( |S_h(x)| \).

We successively compute:

- If \( S_h(x) > 0 \),
  \[
  R_h(x) = S_h(x) - C_u(k_x), \\
  R_m(x) = S_m(x) - C_m(k_x), \\
  R_o(x) = S_o(x) - C_o(k_x).
  \]

- Else,
  \[
  R_h(x) = S_h(x) + C_u(k_x), \\
  R_m(x) = S_m(x) + C_m(k_x), \\
  R_o(x) = S_o(x) + C_o(k_x).
  \]

Again, \( R_h(x) \) and \( R_m(x) \) are exactly representable (hence, they are exactly computed) in double-precision arithmetic:

- \( R_h(x) \) has an absolute value less than \( \pi/4 \) and is a multiple of \( 2^{-49} \),
- \( R_m(x) \) has an absolute value less than \( 2^{-47} + 2^{-50} \) and is a multiple of \( 2^{-99} \).

\( |R_o(x)| \) is less than \( 2^{-97} + 2^{-100} \) and it is computed with error less than or equal to \( 2^{-149} + 2^{-150} + 2^{-154} = 49 \times 2^{-154} \).

- \( 2^{-149} \) is the error bound on \( S_h(x) \);
- \( 2^{-154} \) bounds the error due to the floating-point representation of \( C_o(k_x) \);
- \( 2^{-150} \) bounds the rounding error that occurs when computing \( S_h(x) \pm C_o(k_x) \) in round-to-nearest mode.

Therefore, the number \( R(x) = R_h(x) + R_m(x) + R_o(x) \) is equal to \( x \) minus an integer multiple of \( \pi/2 \) plus an error bounded by \( 49 \times 2^{-154} < 2^{-148} \).

This step is also used (alone, without the previous steps) to reduce small input arguments, less than 8. This allows
our algorithm to perform range-reduction for both kind of arguments, small and medium size. The reduced argument is now stored as the sum of three double-precision numbers, \( R_{hi}(x), R_{med}(x), \) and \( R_{lo}(x) \). We want to return the reduced argument as the sum of two double-precision numbers (one double-precision number may not suffice if we wish to compute trigonometric functions with very good accuracy). To do that, we will use the Fast2sum algorithm presented hereafter.

### 2.3 Final Step

We will get the final result of the range-reduction as follows: Let \( p \) be an integer parameter, \( 1 \leq p \leq 44 \), used to specify the required accuracy. This choice comes from the fact that we work in double precision arithmetic and that, in the most frequent cases, the final relative error will be bounded by \( 2^{-100+p} \). To allow an accurate double precision function result even in the very worst case, we must have a relative error significantly less than \( 2^{-53} \). The problem here is only to propagate the possible carry when summing the three components \( R_{hi}(x), R_{med}(x), \) and \( R_{lo}(x) \). This is performed using floating-point addition and the following result:

**Theorem 2 (Fast2sum algorithm) [7, p. 221, Theorem C].** Let \( a \) and \( b \) be floating-point numbers, with \( |a| \geq |b| \). Assume the used floating-point arithmetic provides correctly rounded results with rounding to the nearest. The following algorithm:

\[
\text{fast2sum}(a, b) : \\
\quad s := a + b \\
\quad z := s - a \\
\quad r := b - z
\]

computes two floating-point numbers \( s \) and \( r \) that satisfy:

- \( r + s = a + b \) exactly;
- \( s \) is the floating-point number which is closest to \( a + b \).

We now consider the different possible cases:

- If \( |R_{hi}(x)| > 1/2^p \), then, since \( |R_{med}(x)| < 2^{-47} + 2^{-50} \), the reduced argument will be close to \( R_{hi}(x) \). In that case, we first compute

\[
t_{med}(x) = R_{med}(x) + R_{lo}(x).
\]

The error on \( t_{med}(x) \) is bounded by the former error on \( R_{lo}(x) \) plus the rounding error due to the addition. Assuming rounding to nearest, this last error is less than or equal to \( 2^{-100} \). Hence, the error on \( t_{med}(x) \) is less than or equal to \( 2^{-100} + 2^{-148} \). Then, we perform (without rounding error)

\[
(y_{hi}, y_{lo}) = \text{fast2sum}(R_{hi}(x), t_{med}(x)).
\]

After that, the two floating-point numbers \( (y_{hi}, y_{lo}) \) represent the reduced argument with an absolute error bounded by \( 2^{-100} + 2^{-148} \approx 2^{-100} \). Hence, the relative error on the reduced argument will be bounded by a value very close to \( 2^{-100+p} \).

- If \( R_{hi}(x) = 0 \), then we perform

\[
(y_{hi}, y_{lo}) = \text{fast2sum}(R_{med}(x), R_{lo}(x)).
\]

After that, since the absolute value of the reduced argument is always larger than \( 0.71 \times 2^{-61} \), the two floating-point numbers \( (y_{hi}, y_{lo}) \) represent the reduced argument with a relative error smaller than

\[
\frac{49 \times 2^{-154}}{0.71 \times 2^{-61}} < 2^{-86}.
\]

- If \( 0 < |R_{hi}(x)| \leq 2^{-p} \), then, since the absolute value of the reduced argument is always larger than \( 0.71 \times 2^{-61} \) and since \( |R_{hi}(x)| < 2^{-97} + 2^{-100} \), most of the information on the reduced argument is in \( R_{hi}(x) \) and \( R_{med}(x) \). We first perform

\[
(y_{hi}, t_{med}) = \text{fast2sum}(R_{hi}(x), R_{med}(x)).
\]

Let \( k \) be the integer satisfying

\[
2^{-k} \leq |y_{hi}| < 2^{-k+1}.
\]

We easily find

\[
|t_{med}| \leq 2^{-k-53}.
\]

After that, we compute

\[
y_{lo} = t_{med} + R_{lo}(x).
\]

The rounding error due to this addition is bounded by \( 2^{-k-107} \). Hence, the two floating-point numbers \( (y_{hi}, y_{lo}) \) represent the reduced argument with an absolute error smaller than

\[
49 \times 2^{-154} + \max\{2^{-107}, 2^{-150}\}.
\]

Therefore, \( (y_{hi}, y_{lo}) \) represent the reduced argument with a relative error better than

\[
49 \times 2^{-154+k} + \max\{2^{-107}, 2^{-150+k}\},
\]

which is less than \( 2^{-87} \) since the absolute value of the reduced argument is less than \( 0.71 \times 2^{-61} \), which implies \( 2^{-k} < 2^{-61} \).

A first solution is to try to make the various error bounds equal. This is done by choosing \( p = 14 \). By doing that, in the worst case, the bound on the relative error will be \( 2^{-86} \), which is quite good. We should notice that, in this case, assuming (10) with \( C = \pi/2 \), the probability that \( |R_{hi}(x)| \) will be less than \( 2^{-p} \) is around \( 7.8 \times 10^{-5} \).

A possibly better solution is to make the most frequent case (i.e., \(|R_{hi}(x)| > 2^{-p}\)) more accurate and to assume that a more accurate yet slower algorithm is used in the other cases (an easy solution is to split the variables into four floating-point values instead of three as we did here). This is done by using a somewhat smaller value of \( p \). For instance, with \( p = 10 \) and \( C = \pi/2 \), still assuming (10), the probability that \( |R_{hi}(x)| < 2^{-p} \) is around \( 1.25 \times 10^{-3} \). In the most frequent case (\(|R_{hi}(x)| \geq 2^{-p}\)), the error bound on the computed reduced argument will be \( 2^{-90} \). Due to its low probability, the other case can be processed with an algorithm a hundred times slower without significantly changing the average time of computation, cf. Amdahl’s law.
2.4 The Algorithm
We can now sketch the complete algorithm:

Algorithm Range-Reduction:

Input: A double-precision floating-point number $x > 0$ and an integer $p > 0$ specifying the required precision in bits.
Output: The reduced argument $y$ given as the sum of two double-precision floating-point numbers $y_{hi}$ and $y_{lo}$, such that $-\pi/4 \leq y < \pi/4$ and $x = k \pi/4 + y$ within an error given in the analysis of Section 2.3, for some integer $k$.

Method:

if $x \geq 2^p - 1$ then
  [Apply the method of Payne and Hanek.]
else if $x \leq 8$ then
  \[ S_{hi} \leftarrow x; S_{med} \leftarrow 0; S_{lo} \leftarrow 0; \]
else
  \[ I \leftarrow \text{round}(x); \rho \leftarrow x - I; \]
  \[ S_{hi} \leftarrow \rho; S_{med} \leftarrow 0; S_{lo} \leftarrow 0; \]
  \[ i \leftarrow 7; j \leftarrow 56; \]
  while $i \geq 0$ do
    \[ w \leftarrow \text{round}(I >> j); \]
    \[ S_{hi} \leftarrow S_{hi} + \text{sign}(w)T_{hi}(i, |w|); \]
    \[ S_{med} \leftarrow S_{med} + \text{sign}(w)T_{med}(i, |w|); \]
    \[ i \leftarrow i - 1; j \leftarrow j - 8; \]
  \[ S_{lo} \leftarrow \sum_i^j \text{sign}(w)T_{lo}(i, |w|) \text{ (cf. 11)}; \]
if \[ |S_{hi}| \geq \pi/4 \text{ then} \]
  \[ k \leftarrow \text{Reduce}(|S_{hi}|); \]
  \[ S_{hi} \leftarrow S_{hi} + \text{sign}(S_{hi})C_{hi}(k); \]
  \[ S_{med} \leftarrow S_{med} + \text{sign}(S_{hi})C_{med}(k); \]
  \[ S_{lo} \leftarrow S_{lo} + \text{sign}(S_{hi})C_{lo}(k); \]
else if \[ |S_{hi}| > 2^p \text{ then} \]
  \[ \text{temp} \leftarrow S_{med} + S_{lo}; \]
  \[ (y_{hi}, y_{lo}) \leftarrow \text{fast2sum}(S_{hi}, \text{temp}); \]
else if \[ S_{hi} = 0 \text{ then} \]
  \[ (y_{hi}, y_{lo}) \leftarrow \text{fast2sum}(S_{med}, S_{lo}); \]
else \[ (y_{hi}, \text{temp}) \leftarrow \text{fast2sum}(S_{hi}, S_{lo}); \]
  \[ y_{lo} \leftarrow \text{temp} + S_{lo}; \]

Where: The function \text{Reduce}(|S_{hi}|) chooses the appropriate multiple $k$ of $\pi/2$, represented as the triple $(C_{hi}(k), C_{med}(k), C_{lo}(k))$.

3 Cost of the Algorithm
In this section, we compare our method to other algorithms on the same input range $[2^{2^{30}} - 1]$; Payne and Hanek’s methods (see Section 1.1) and the Modular range-reduction method described in [1]. Concerning Payne and Hanek’s method, we used the version of the algorithm used by Sun Microsystems [10]. We chose as criteria for the evaluation of the algorithms the table size, the number of table accesses, and the number of floating-point multiplications, divisions, and additions.

Table 2 shows the potential advantages of our algorithm for small and medium-sized input argument. Payne and Hanek’s method over that range doesn’t need much memory, but roughly requires three times as many operations. The Modular range-reduction has the same characteristics as Payne and Hanek’s method concerning the table size needed and the number of elementary operations involved, but requires more table accesses. Our algorithm is then a good compromise between table size and number of operations for range-reduction of medium-sized argument.

To get more accurate figures than by just counting the operations, we have implemented this algorithm in ANSI-C. The program can be downloaded from http://gala.univ-perp.fr/~ddefour/high_radix.tgz. This implementation shows that our algorithm is 4 to 5 times faster, depending on the required final precision, than the Sun implementation of Payne and Hanek’s algorithm, provided that the tables are in main memory (which will be true when the trigonometric functions are frequently called in a numerical program; and, when they are not frequently called, the speed of range-reduction is no longer an issue). Our algorithm is then a good compromise between table size and delay for range-reduction of small and medium-sized arguments.

A variant of our algorithm would consist of first computing $S_{hi}, S_{med}$ and $R_{hi}, R_{med}$ only. Then, during the fourth step of the algorithm, if the accuracy does not suffice, compute $T_{lo}$ and $R_{lo}$. This slight modification can reduce the number of elementary operations in the (most frequent) cases where no extra accuracy is needed. We can also reduce the table size by 4 Kbytes by storing the $T_{lo}$ values in single-precision only, instead of using double-precision.

Another variant (that can be useful depending on the processor and compiler) would be to replace the loop “while $i \geq 0$” with “while $I < 0$ and $i \geq 0.$” In that case (for a medium-sized argument $x$), the number $N$ of double-precision floating-point operations becomes at most

\[ N = 17 + 2[\log_{256} x], \]

i.e., $19 \leq N \leq 33$. Also, the number of table accesses becomes at most $11 + 2[\log_{256} x]$.

4 Conclusions
We have presented an algorithm for accurate range-reduction of input arguments with absolute value less than $2^{30} - 1$. This table-based algorithm gives accurate results for the most frequent cases. In order to cover the whole double-precision domain for input arguments, we suggest using Payne and Hanek’s algorithm for huge arguments. A major
drawback of our method lies in the table size needed, thus a future effort will be to reduce the table size, while keeping a good trade off between speed and accuracy.

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